

How to find dead ends efficiently

A node is a dead end if it has no out-link, or all of its out-links point to dead ends only. To find dead end efficiently (in $O(n + m)$ time), you should use two arrays (of lists):

- Array N^+ , where $N^+[i]$ is a *list* of nodes that *links to* i . In other words, $N^+[i]$ is the set of in-neighbors of i . Note that $len(N^+[i])$ is the number of in-links of i .
- Array N^- , where $N^-[i]$ is a *list* of nodes that i *links to*. In other words, $N^-[i]$ is the set of out-neighbors of i . Note that $len(N^-[i])$ is the number of out-links of i .

In addition, you should use an array D to store the out-degree of i . That is, $D[i] = leng(N^-[i])$. Dead ends would have $D[i] = 0$. Also, you should use a Queue, say q (a FIFO queue) to store temporary dead-ends during the execution of the algorithm below:

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FINDDEADENDS( $G(V, E), N^+, N^-$ )
  initialize array  $D$  so that  $D[i] = 0$  for all node  $i$ 
  initialize an empty Queue  $q$ .
  for each node  $i \in V$ 
     $D[i] = len(N^-[i])$ 
    if  $D[i] = 0$ 
      put  $i$  to  $q$  //  $i$  is a dead end
  initialize an empty list  $L$ 
  while  $q$  is not empty
     $i \leftarrow$  pop an element from  $q$ 
    if  $i$  is not in  $L$ 
      append  $i$  to the end of  $L$ 
      for each  $j$  in  $N^+[i]$ 
         $D[j] \leftarrow D[j] - 1$  // remove  $i$  will decrease out-deg of  $j$ 
        if  $D[j] = 0$  //  $j$  is a new dead end
          put  $j$  to Queue  $q$ 
  return  $L$ 
```

The return list L would be the set of dead ends. Their order in L is their removal order.

Tip: to check whether a node i in L , you just need to use a boolean array, say M . Initially, every node i has $M[i] = \text{False}$. Every time you put a node i to L , mark $M[i] = \text{True}$.